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REVERSE ENGINEERING GOLANG BINARIES WITH GHIDRA

Dorka Palotay & Albert Zsigovits

CUJO AI, Hungary

dorka.palotay@cujo.com

albert.zsigovits@cujo.com

INTRODUCTION

Go (also called Golang) is an open-source programming language that was designed by *Google* in 2007 and made available to the public in 2012. Over the years it has gained popularity among developers and, as usually happens, it has not only become popular with developers of legitimate software but has also attracted the attention of malware developers. The fact that Go supports cross compiling to run binaries on various operating systems makes it a tempting choice for malware developers. The possibility to compile the same code for all major platforms (*Windows*, *Linux* and *MacOS*) makes the attackers' lives much easier, as they don't have to develop and maintain a different codebase for each target environment.

Some special features of the Go programming language make investigating Go binaries difficult for reverse engineers. Reverse engineering tools (e.g. disassemblers) can do a great job in analysing binaries that are written in more popular languages (e.g. C, C++, .NET), but Go presents new challenges that makes the analysis more cumbersome.

Go binaries are usually statically linked, which means that all the necessary libraries are included in the compiled binary. This results in large binaries. On the one hand this makes malware distribution more difficult for the attackers, but on the other hand some security products also have issues with handling such large files. The other advantage of statically linked binaries for the attackers is that the malware can run on the target systems without dependency issues.

As we see a continuous growth in malware written in Go, and we expect more families to emerge, we decided to dive deeper into the Go programming language and enhance our toolset to be more effective in investigating Go malware.

In the first section of this paper we provide a list of the recently discovered malware families written in Go and briefly introduce a few of them.

In the next sections we will discuss two of the difficulties that reverse engineers face during Go binary analysis and we will show our solutions for those.

Ghidra [1] is an open-source reverse engineering tool developed by the National Security Agency, which we frequently use for static malware analysis. It is possible to create custom scripts and plug-ins for Ghidra to provide specific functionalities that are needed by researchers. We used this feature of Ghidra and created custom scripts to aid our Go binary analysis.

In our research we tested Go until version 1.15 and used Ghidra versions 9.1 and 9.2.3.

The slides and other materials accompanying this paper are available in our *GitHub* repository [2].

GO MALWARE FAMILIES

In this section, we will briefly look at some of the prominent Go malware families. Table 1 shows a list of the recently discovered malware families written in Go, some of which we introduce in the following sections.

FritzFrog P2P botnet

This piece of malware was discovered by *Guardicore* [27]. FritzFrog has been active since January 2020. With its decentralized nature, there is no single command-and-control server, which makes it very unique as a Peer-2-Peer (P2P) botnet. Its worm executable is completely written in Golang, and its P2P implementation is proprietary.

FritzFrog is also considered a highly advanced piece of malware due to its multi-threaded, modular and fileless nature, which is very rare in a Mirai- and Gafgyt-variant dominated world.

Once a victim is successfully breached, it starts running the UPX-packed malware, which immediately erases itself. The malware process runs under the names *ifconfig* and *nginx*, to minimize suspicion.

Its main targets were governmental offices, educational institutions, medical centres, banks and numerous telecom companies as it tried to infiltrate via brute-force through the SSH protocol.

Guardicore has also found FritzFrog to have some similarity to the Rakos botnet, as its function naming is similarly written, and its version numbers are very much alike. They have also developed a client program, which can send commands to the botnet by injecting its own node to participate in the P2P network.

The final goal of the malware is to deploy the malicious payload of a Monero cryptocurrency miner. FritzFrog has been observed with 20 different versions and variants since its inception.

HEH P2P botnet

Another botnet that made headlines as Go malware is the HEH botnet, discovered by *360 Netlab* [16].

HEH's initial vector of attack is the Telnet protocol, on port 23 or 2323, by brute-forcing its way through the login prompt. In the analysed variants, there were 171 usernames and 504 potential passwords stored in variables.

HEH uses a proprietary P2P protocol. HEH also has three clear distinct modules: a propagation module, an HTTP service module and a P2P module.

Family	Reference
Kaiji	Intezer - New Chinese Linux malware using Golang [3]
Zebrocy	A Zebrocy Go Downloader [4]
eCh0riax	Reverse Engineering Go Binaries with Ghidra - CUJO AI [5]
LiquorBot	Intezer on Twitter [6]
WellMess	Intezer on Twitter [7]
Smaug ransomware	Anomali Threat Research Releases First Public Analysis of Smaug Ransomware as a Service [8]
FritzFrog	FritzFrog: A New Generation Of Peer-To-Peer Botnets - Guardicore [9]
Godlike12	Holy water: ongoing targeted water-holing attack in Asia [10]
IRCFly	muesli/ircfly [11]
IPStorm	The InterPlanetary Storm: New Malware in Wild Using InterPlanetary File System's (IPFS) p2p network [12]
Nephilim	Vitali Kremez on Twitter [13]
EKANS	EKANS Ransomware Targeting OT ICS Systems FortiGuard Labs [14]
RobinHood	Vitali Kremez on Twitter [15]
HEH	https://blog.netlab.360.com/heh-an-iot-p2p-botnet/ [16]
Go Loader	TA416 Goes to Ground and Returns with a Golang PlugX Malware Loader Proofpoint US [17]
GOSH	Intezer on Twitter [18]
Glupteba.Go	Glupteba malware hides in plain sight [19]
New RAT	There's a New a Golang-written RAT in Town [20]
BlackRota	https://blog.netlab.360.com/blackrota-an-obfuscated-backdoor-written-in-go-en/ [21]
Clipboard.Stealer	https://analyze.intezer.com/files/bd978ba0d723aea3106c6abc58cf71df5abe4d674d0d1fc38b37d4926d740738 [22]
CryptoStealer.Go	Analyzing a new stealer written in Golang - Malwarebytes Labs [23]
Sysrv-hello	Sysrv Botnet Expands and Gains Persistence Official Juniper Networks Blogs [24]
Epsilon Red	A new ransomware enters the fray: Epsilon Red [25]
aicm	Intezer on Twitter [26]

Table 1: Go malware families.

According to 360 Netlab, this botnet is not yet mature, as some of the more essential functions, like the attack module, have not been implemented yet, and there are flaws in the implementation of the P2P module too.

HEH starts with a shell script, which pulls down the malicious binaries to different types of architectures and, surprisingly, executes all of them on the target. The malicious binary then kills a series of service processes based on listening port numbers.

HEH also starts an HTTP server on TCP port 80, and an initial dummy content will be placed onto the server, which gets overwritten by the P2P module once data is transferred from another node.

Currently, the botnet can execute shell commands, update the Peer List and exchange data, but as the attack module is not yet finished, analysts expect that there will be several iterations of HEH versions.

Sysrv botnet

In April, researchers at *Juniper Threat Labs* [24] reported that they had discovered a surge of activity from the botnet Sysrv. Traces of Sysrv botnet activity date back to December 2020.

Previously, Sysrv had separate worm and miner executables, but more recently Sysrv combines the two in one malicious binary. We also know that Sysrv once used two mining pools but now focuses only on the miner pool, 'nanopool'.

Some developments have been observed in the loader script itself, which loads the malicious binary: the script now involves a procedure for adding an SSH key to the `authorized_keys` file on the target system to achieve persistence. Also, there is a *Linux* version of loader script, which is called `ldr.sh`, and a *Windows* one, called `ldr.ps1`.

The first variants of the malicious Sysrv payload exploited several different vulnerabilities, including the following:

- CVE-2020-16846 – Saltstack RCE
- CVE-2019-10758 – Mongo Express RCE
- CVE-2018-7600 – Drupal Ajax RCE
- CVE-2017-11610 – XML-RPC
- XXL-JOB Unauth RCE (without CVE)
- ThinkPHP RCE

Later versions of Sysrv started to include many other application-specific exploits, and we expect that they will keep incorporating more. These application-specific exploits are used to download and execute the first-stage loader script, `ldr.sh` or `ldr.ps1`.

Sysrv's goal is to spread further and deploy a Monero cryptocurrency miner on the infected systems.

Epsilon Red ransomware

Researchers at *Sophos* [25] discovered a Golang-based ransomware that was attacking a US-based business. The loader for the ransomware payload is a PowerShell script.

Analysts conclude that this new ransomware variant is quite a simple program, as it has no networking capabilities, and the encryption process is simple.

Epsilon Red will encrypt everything in its way, including all system files, possibly rendering the entire operating system unusable. Once the encryption is done, the ransomware appends the extension '.epsilonired' to all encrypted files. The ransomware spawns a new child process for every folder it encrypts, which results in an unnecessarily long list of running ransomware processes.

From a binary perspective, the malicious sample was compiled with MinGW, and packed with a modified version of the UPX packer. We have also observed that the sample contains code from the open-source project Godirwalk: this tool will scan the entire system storage and compile a list of directory paths, which is then used for the encryption.

Analysts have found that the ransomware note dropped by Epsilon has some similarity to the one left behind by the REvil ransomware.

We have made the following observations and predictions during the analysis of the aforementioned botnets:

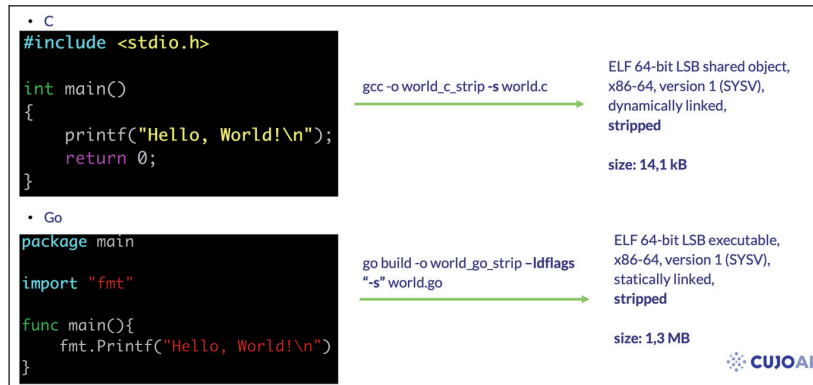
- Ransomware that is written in Golang will become more common
- P2P botnets are still popular and introduce new concepts and modules
- Botnets still trying to deploy cryptocurrency miners as a final step

Due to these, we have decided to dive deep into the Go language to understand it better and to enhance our ability to tackle Go malware. In the next two sections we introduce two features of Go, the difficulties reverse engineers face during Go malware analysis thanks to those, and our solutions.

LOST FUNCTION NAMES

The first issue is not specific to Go binaries, but stripped binaries in general. Compiled executable files can contain debug symbols which make debugging and analysis easier. When reverse engineering a program that was compiled with debugging information included, analysts can see not only memory addresses but also the names of the routines and variables. However, in order to reduce the size, developers usually compile the files without this information, creating so-called stripped binaries. For malware authors another advantage of stripping binaries is that it makes reverse engineering more difficult. In this case analysts cannot rely on the function names to help them find their way around the code. For statically linked Go binaries, where all the necessary libraries are included, this can significantly slow down the analysis.

To illustrate this issue, we used simple 'Hello World' examples written in C⁽¹⁾ and Go⁽²⁾ for comparison and compiled them to stripped binaries. Note the size difference between the two executables.

Figure 1: Hello World examples written in C⁽¹⁾ and Go⁽²⁾.

Ghidra's function window lists all the defined functions within the binaries. In the non-stripped versions, function names are nicely visible and provide a great help for reverse engineers.

Figure 2 shows the function list for the world_c binary. The table lists 19 items.

Name	Location	Function Signature	Function Size
_init	00101000	int _init(EVP...	27
FUN_00101020	00101020	undefined FUN...	13
_cxa_finalize	00101040	thunk undefined...	11
puts	00101050	thunk int puts...	11
_start	00101060	undefined _sta...	47
deregister_tm_clones	00101090	undefined dere...	34
register_tm_clones	001010c0	undefined regi...	51
_do_global_ctors_aux	00101100	undefined _do...	54
frame_dummy	00101140	thunk undefined...	9
main	00101149	undefined main()	27
__libc_csu_init	00101170	undefined __li...	101
__libc_csu_fini	001011e0	undefined __li...	5
_fini	001011e8	undefined _fini()	13
_ITM_deregisterTMCloneTable	00105000	thunk undefined...	1
puts	00105008	thunk int puts...	1
__libc_start_main	00105010	thunk undefined...	1
_gmon_start_	00105018	thunk undefined...	1
_ITM_registerTMCloneTable	00105020	thunk undefined...	1
_cxa_finalize	00105028	thunk undefined...	1

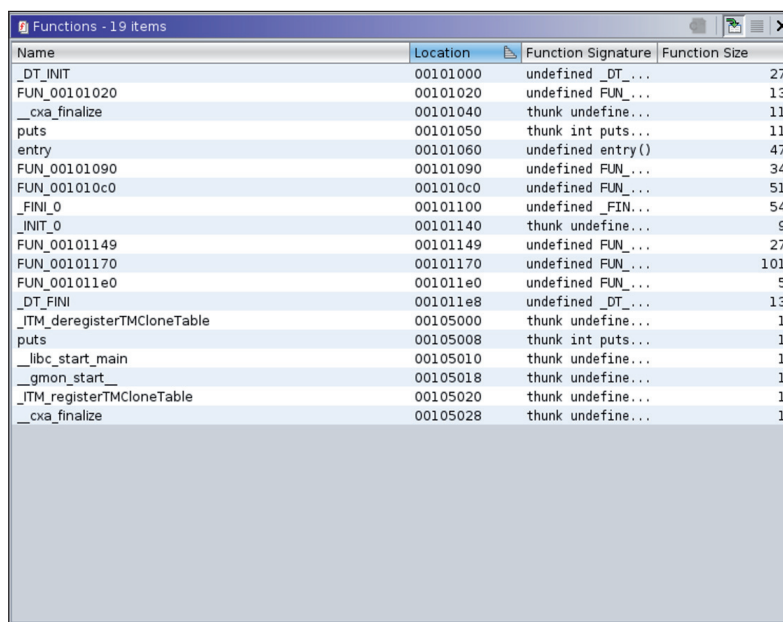
Figure 2: world_c⁽³⁾ function list.

Figure 3 shows the function list for the world_go binary. The table lists 1790 items.

Name	Location	Function Signat...	Function Size
internal/cpu.initialize	00401000	undefined int...	78
internal/cpu.processOptions	00401060	undefined int...	1877
internal/cpu.indexByte	004017c0	undefined int...	53
internal/cpu.doinit	00401800	undefined int...	1029
internal/cpu.cpuid	00401c20	undefined int...	27
internal/cpu.xgetbv	00401c40	undefined int...	17
type..eq.internal/cpu.CacheLinePad	00401c60	undefined typ...	6
type..eq.internal/cpu.option	00401c80	undefined typ...	165
type..eq.[15]internal/cpu.option	00401d40	undefined typ...	139
runtime/internal/sys.OnesCount64	00401de0	undefined run...	119
runtime/internal/atomic.Cas64	00401e60	undefined run...	26
runtime/internal/atomic.Casuintptr	00401e80	thunk undefin...	5
runtime/internal/atomic.Storeuintptr	00401ea0	thunk undefin...	5
runtime/internal/atomic.Store	00401ec0	undefined run...	12
runtime/internal/atomic.Store64	00401ee0	undefined run...	14
internal/bytealg.init.0	00401f00	undefined int...	34
cmpbody	00401f40	undefined cmp...	569
runtime.cmpstring	00402180	undefined run...	30
memeqbody	004021a0	undefined mem...	318
runtime.memequal	004022e0	undefined run...	36
runtime.memequal_varlen	00402320	undefined run...	35
indexbytebody	00402360	undefined ind...	279
internal/bytealg.IndexByteString	00402480	undefined int...	24
runtime.memhash128	004024a0	undefined run...	89
runtime.strhashFallback	00402500	undefined run...	99
runtime.f32hash	00402580	undefined run...	282
runtime.f64hash	004026a0	undefined run...	284
runtime.c64hash	004027c0	undefined run...	110
runtime.c128hash	00402840	undefined run...	110

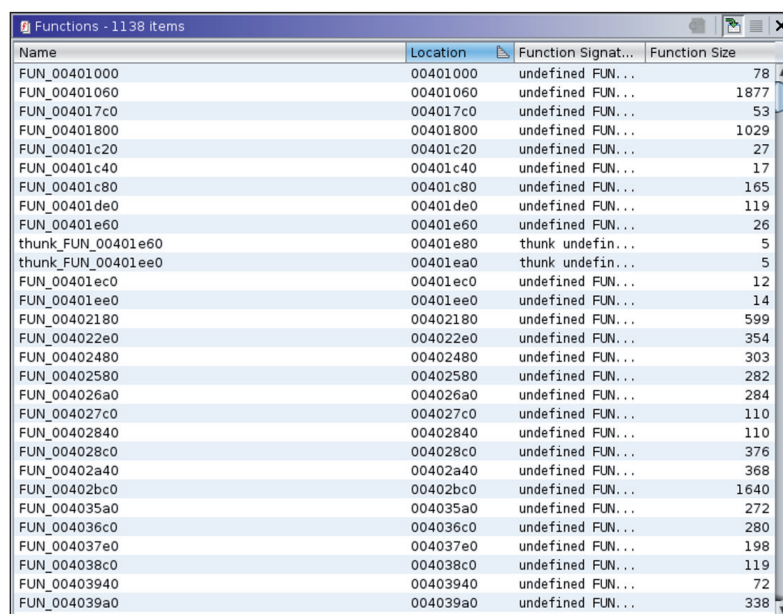
Figure 3: world_go⁽⁵⁾ function list.

For stripped binaries the function lists look the following:



Name	Location	Function Signature	Function Size
DT_INIT	00101000	undefined DT...	27
FUN_00101020	00101020	undefined FUN...	13
_cxa_finalize	00101040	thunk undefined...	11
puts	00101050	thunk int puts...	11
entry	00101060	undefined entry()	47
FUN_00101090	00101090	undefined FUN...	34
FUN_001010c0	001010c0	undefined FUN...	51
_FINI_0	00101100	undefined FIN...	54
_INIT_0	00101140	thunk undefined...	9
FUN_00101149	00101149	undefined FUN...	27
FUN_00101170	00101170	undefined FUN...	101
FUN_001011e0	001011e0	undefined FUN...	5
_DT_FINI	001011e8	undefined DT...	13
_ITM_deregisterTMCloneTable	00105000	thunk undefined...	1
puts	00105008	thunk int puts...	1
_libc_start_main	00105010	thunk undefined...	1
_gmon_start	00105018	thunk undefined...	1
_ITM_registerTMCloneTable	00105020	thunk undefined...	1
_cxa_finalize	00105028	thunk undefined...	1

Figure 4: *world_c_strip*⁽⁴⁾ function list.



Name	Location	Function Signature	Function Size
FUN_00401000	00401000	undefined FUN...	78
FUN_00401060	00401060	undefined FUN...	1877
FUN_004017c0	004017c0	undefined FUN...	53
FUN_00401800	00401800	undefined FUN...	1029
FUN_00401c20	00401c20	undefined FUN...	27
FUN_00401c40	00401c40	undefined FUN...	17
FUN_00401c80	00401c80	undefined FUN...	165
FUN_00401de0	00401de0	undefined FUN...	119
FUN_00401e60	00401e60	undefined FUN...	26
thunk_FUN_00401e60	00401e80	thunk undefin...	5
thunk_FUN_00401ee0	00401ea0	thunk undefin...	5
FUN_00401ec0	00401ec0	undefined FUN...	12
FUN_00401ee0	00401ee0	undefined FUN...	14
FUN_00402180	00402180	undefined FUN...	599
FUN_004022e0	004022e0	undefined FUN...	354
FUN_00402480	00402480	undefined FUN...	303
FUN_00402580	00402580	undefined FUN...	282
FUN_004026a0	004026a0	undefined FUN...	284
FUN_004027c0	004027c0	undefined FUN...	110
FUN_00402840	00402840	undefined FUN...	110
FUN_004028c0	004028c0	undefined FUN...	376
FUN_00402a40	00402a40	undefined FUN...	368
FUN_00402bc0	00402bc0	undefined FUN...	1640
FUN_004035a0	004035a0	undefined FUN...	272
FUN_004036c0	004036c0	undefined FUN...	280
FUN_004037e0	004037e0	undefined FUN...	198
FUN_004038c0	004038c0	undefined FUN...	119
FUN_00403940	00403940	undefined FUN...	72
FUN_004039a0	004039a0	undefined FUN...	338

Figure 5: *world_go_strip*⁽⁶⁾ function list.

These examples show nicely that even a simple ‘hello world’ Go binary is huge, with more than a thousand functions, and in the stripped version reverse engineers cannot rely on the function names to aid their analysis.

Note: As a result of stripping, not only did the function names disappear, but instead of 1,790 defined functions only 1,138 were recognized by Ghidra.

We were interested to find out whether there is a way to recover the function names within stripped binaries. First, using a simple string search we can check if the function names are still available within the binaries. For the C example we were looking for the function name ‘main’, while in the Go example it is ‘main.main’.

```
> strings world_c | grep -o ".\{0,10\}main.\{0,10\}"
ibc_start_main
ibc_start_main@@GLIBC_2.
main
```

Figure 6: *world_c*⁽³⁾ strings – ‘main’ was found.

```
> strings world_c_strip | grep -o ".\{0,10\}main.\{0,10\}"
ibc_start_main
```

Figure 7: world_c_strip⁽⁴⁾ strings – ‘main’ was not found.

```
> strings world_go | grep -o ".\{0,10\}main.\{0,10\}"
hasmain
edruntime.main not on m0
p stateremaining pointe
out of domainpanic whil
e space remainingreflect
routines (main called ru
runtime.main
runtime.main.func1
runtime.main.func2
main.main
main..inittask
runtime.main_init_done
runtime.mainStarted
runtime.mainPC
runtime.main
runtime.main.func1
runtime.main.func2
main.main
```

Figure 8: world_go⁽⁵⁾ strings – ‘main.main’ was found.

```
> strings world_go_strip | grep -o ".\{0,10\}main.\{0,10\}"
hasmain
edruntime.main not on m0
p stateremaining pointe
out of domainpanic whil
e space remainingreflect
routines (main called ru
runtime.main
runtime.main.func1
runtime.main.func2
main.main
```

Figure 9: world_go_strip⁽⁶⁾ strings – ‘main.main’ was found.

While in the stripped C binary⁽⁴⁾ the function name cannot be found with the strings utility, in the Go version⁽⁶⁾ ‘main.main’ is still available. This discovery gave us some hope that function name recovery might be possible in stripped Go binaries.

Loading the binary⁽⁶⁾ to Ghidra and searching for the ‘main.main’ string will show the exact location. As can be seen in Figure 10, the function name string is located within the .gopclntab section.

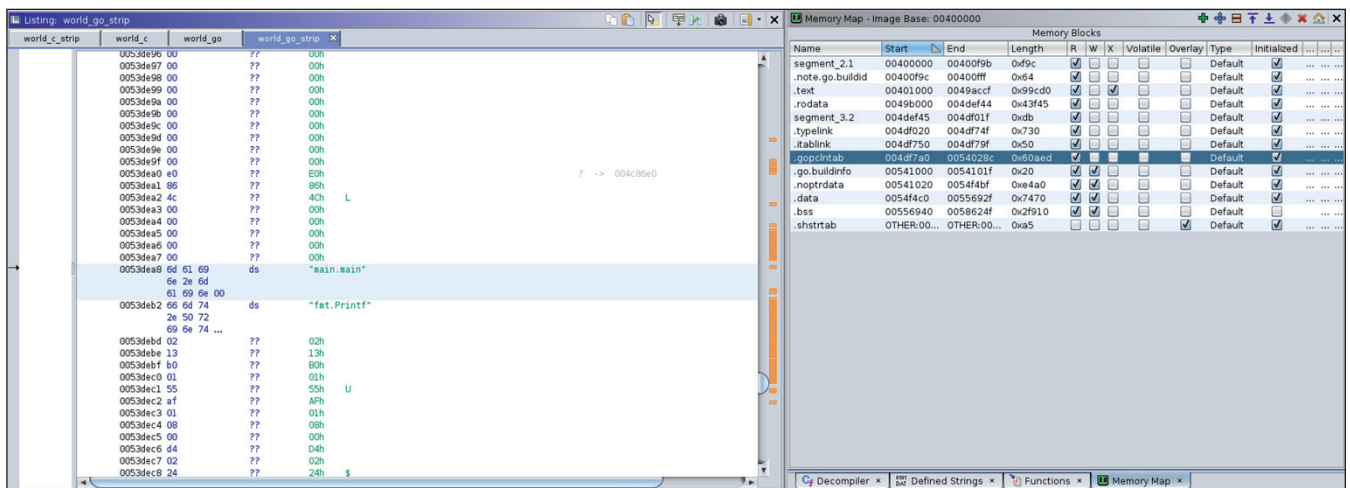


Figure 10: world_go_strip⁽⁶⁾ main.main string in Ghidra.

The pcIntab structure has been available since Go version 1.2 and is nicely documented [28]. The structure starts with a magic value followed by information about the architecture. Then the function symbol table holds information about the functions within the binary. The address of the entry point of each function is followed by a function metadata table.

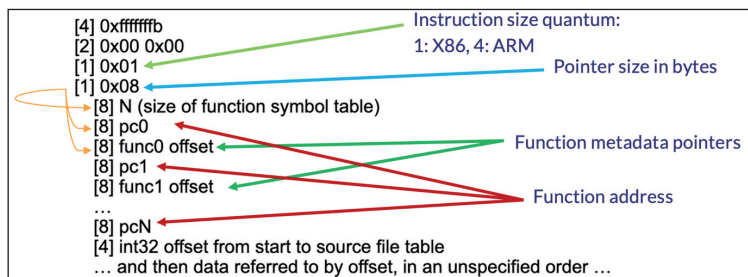


Figure 11: The pcIntab structure.

The function metadata table, among other important information, stores an offset to the function name.

```
struct      Func
{
    uintptr    entry;    // start pc
    int32      name;      // name (offset to C string)
    int32      args;      // size of arguments passed to function
    int32      frame;     // size of function frame, including saved caller PC
    int32      pcsp;       // pcsp table (offset to pcvalue table)
    int32      pcfile;     // pcfile table (offset to pcvalue table)
    int32      pcIn;       // pcIn table (offset to pcvalue table)
    int32      nfuncdata;  // number of entries in funcdata list
    int32      npcdata;    // number of entries in pcdata list
};
```

Figure 12: Function metadata table.

Using this information, it is possible to recover the function names. Our team created a script (go_func.py) for Ghidra to recover function names in stripped Go ELF files by executing the following steps:

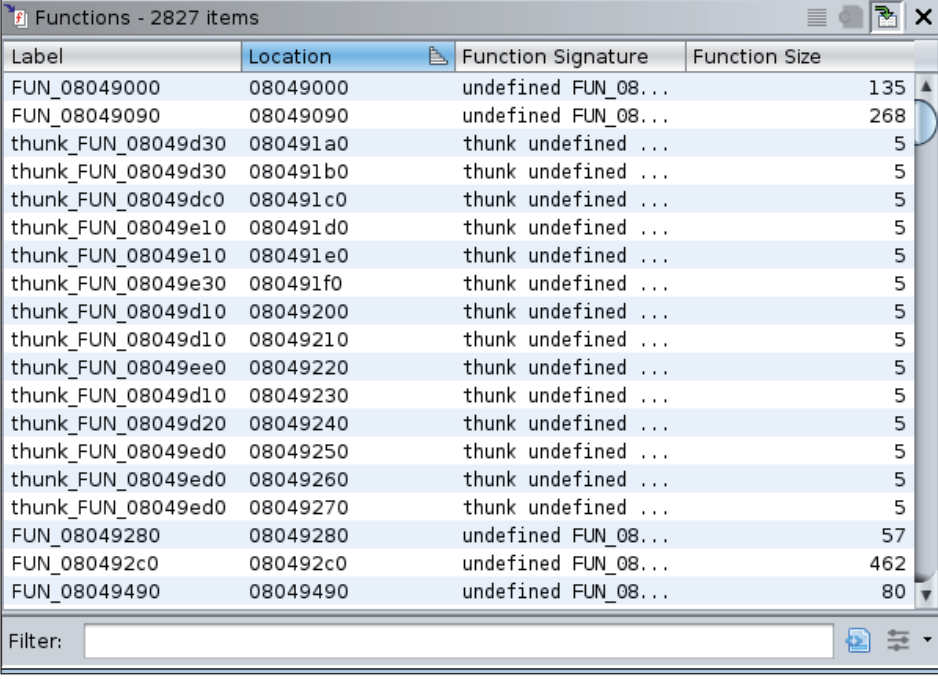
- Locate pcIntab structure
- Extract function addresses
- Find function name offsets

After executing our script not only will the function names be restored, but the previously unrecognized functions will be defined as well.

Functions - 1790 items			
Name	Location	Function Signat...	Function Size
fmt.(*pp).Flag	00492de0	undefined fmt...	143
fmt.(*pp).Write	00492e80	undefined fmt...	271
fmt.Fprintf	00492fa0	undefined fmt...	268
fmt.getField	004930c0	undefined fmt...	183
fmt.parsenum	00493180	undefined fmt...	219
fmt.(*pp).unknownType	00493260	undefined fmt...	784
fmt.(*pp).badVerb	00493580	undefined fmt...	1649
fmt.(*pp).fmtBool	00493c00	undefined fmt...	111
fmt.(*pp).fmt0x64	00493c80	undefined fmt...	149
fmt.(*pp).fmtInteger	00493d20	undefined fmt...	820
fmt.(*pp).fmtFloat	00494060	undefined fmt...	408
fmt.(*pp).fmtComplex	00494200	undefined fmt...	583
fmt.(*pp).fmtString	00494460	undefined fmt...	457
fmt.(*pp).fmtBytes	00494640	undefined fmt...	2303
fmt.(*pp).fmtPointer	00494f40	undefined fmt...	1358
fmt.(*pp).catchPanic	004954a0	undefined fmt...	1534
fmt.(*pp).handleMethods	00495aa0	undefined fmt...	1748
fmt.(*pp).printArg	00496180	undefined fmt...	2348
fmt.(*pp).printValue	00496ae0	undefined fmt...	9767
fmt.intFromArg	00499140	undefined fmt...	529
fmt.parseArgNumber	00499360	undefined fmt...	293
fmt.(*pp).argNumber	004994a0	undefined fmt...	278
fmt.(*pp).badArgNum	004995c0	undefined fmt...	367
fmt.(*pp).missingArg	00499740	undefined fmt...	367
fmt.(*pp).doPrintf	004998c0	undefined fmt...	4490
fmt.glob..func1	0049aa60	undefined fmt...	84
fmt.init	0049aac0	undefined fmt...	197
type..eq.fmt.fmt	0049aba0	undefined typ...	172
main.main	0049ac60	undefined mai...	112

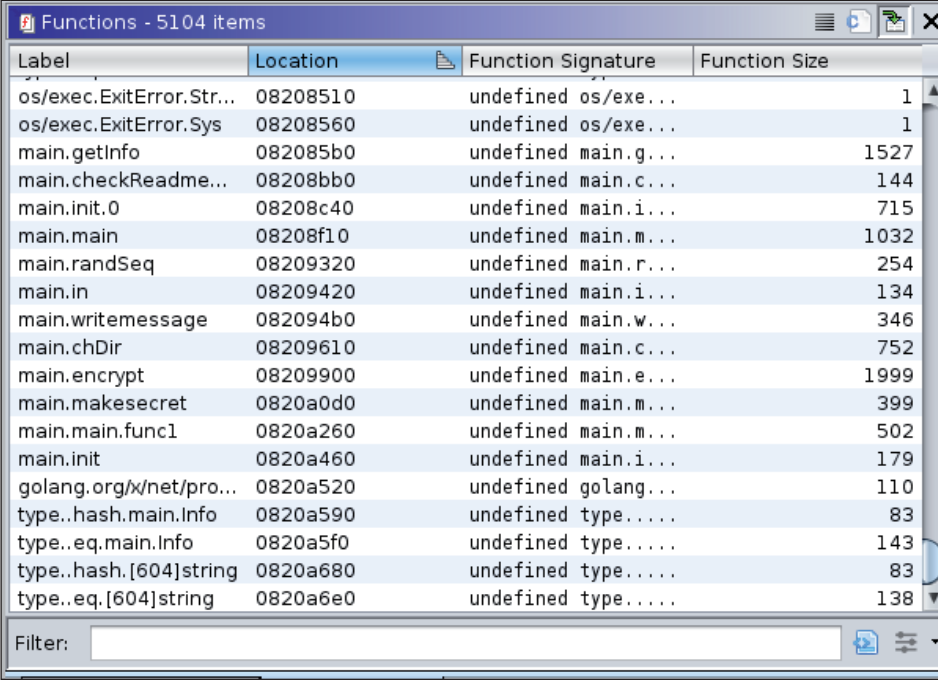
Figure 13: world_go_strip⁽⁶⁾ function list after executing go_func.py.

To see a real-world example let's look at an eCh0raix ransomware sample⁽⁹⁾:



Label	Location	Function Signature	Function Size
FUN_08049000	08049000	undefined FUN_08...	135
FUN_08049090	08049090	undefined FUN_08...	268
thunk_FUN_08049d30	080491a0	thunk undefined ...	5
thunk_FUN_08049d30	080491b0	thunk undefined ...	5
thunk_FUN_08049dc0	080491c0	thunk undefined ...	5
thunk_FUN_08049e10	080491d0	thunk undefined ...	5
thunk_FUN_08049e10	080491e0	thunk undefined ...	5
thunk_FUN_08049e30	080491f0	thunk undefined ...	5
thunk_FUN_08049d10	08049200	thunk undefined ...	5
thunk_FUN_08049d10	08049210	thunk undefined ...	5
thunk_FUN_08049ee0	08049220	thunk undefined ...	5
thunk_FUN_08049d10	08049230	thunk undefined ...	5
thunk_FUN_08049d20	08049240	thunk undefined ...	5
thunk_FUN_08049ed0	08049250	thunk undefined ...	5
thunk_FUN_08049ed0	08049260	thunk undefined ...	5
thunk_FUN_08049ed0	08049270	thunk undefined ...	5
FUN_08049280	08049280	undefined FUN_08...	57
FUN_080492c0	080492c0	undefined FUN_08...	462
FUN_08049490	08049490	undefined FUN_08...	80

Figure 14: eCh0raix⁽⁹⁾ function list.



Label	Location	Function Signature	Function Size
os/exec.ExitError.Str...	08208510	undefined os/exe...	1
os/exec.ExitError.Sys	08208560	undefined os/exe...	1
main.getInfo	082085b0	undefined main.g...	1527
main.checkReadme...	08208bb0	undefined main.c...	144
main.init.0	08208c40	undefined main.i...	715
main.main	08208f10	undefined main.m...	1032
main.randSeq	08209320	undefined main.r...	254
main.in	08209420	undefined main.i...	134
main.writemessage	082094b0	undefined main.w...	346
main.chDir	08209610	undefined main.c...	752
main.encrypt	08209900	undefined main.e...	1999
main.makesecret	0820a0d0	undefined main.m...	399
main.main.func1	0820a260	undefined main.m...	502
main.init	0820a460	undefined main.i...	179
golang.org/x/net/pro...	0820a520	undefined golang...	110
type..hash.main.Info	0820a590	undefined type.....	83
type..eq.main.Info	0820a5f0	undefined type.....	143
type..hash.[604]string	0820a680	undefined type.....	83
type..eq.[604]string	0820a6e0	undefined type.....	138

Figure 15: eCh0raix⁽⁹⁾ function list after executing go_func.py.

This example clearly shows how much help this simple function name recovery script can be during reverse engineering. Only by looking at the function names can analysts assume that they are dealing with a ransomware.

Note: In *Windows* Go binaries there is no specific section for the pcIntab structure, rather researchers need to search explicitly for the fields of this structure (e.g. magic value, possible field values). For *MacOS* the `_gopclntab` section is available, and similarly `.gopclntab` in *Linux* binaries.

Challenges

If a function name string is not defined by Ghidra, then the function name recovery script will fail to rename that specific function, since it cannot find the function name string at the given location. To overcome this issue our script always checks

if a defined data type is located at the function name address and if it isn't, then before renaming a function it tries to define a string data type at the given address.

In the example shown in Figures 16 and 17 the function name string 'log.New' is not defined in an eCh0raix ransomware sample⁽⁹⁾, so the corresponding function cannot be renamed without string creation first.

083aa0e4	6c	??	6Ch	l
083aa0e5	6f	??	6Fh	o
083aa0e6	67	??	67h	g
083aa0e7	2e	??	2Eh	.
083aa0e8	4e	??	4Eh	N
083aa0e9	65	??	65h	e
083aa0ea	77	??	77h	w
083aa0eb	00	??	00h	

Figure 16: eCh0raix⁽⁹⁾ log.New function name undefined.

* FUNCTION *				

undefined FUN_08184fa0(undefined4 param_1, undefined4 pa...				
undefined	AL:1	<RETURN>		
undefined4	Stack[0x4]:4	param_1	XREF[1]:	08184fc7(R)
undefined4	Stack[0x8]:4	param_2	XREF[2]:	08184fd8(R),
				0818501d(R)
undefined4	Stack[0xc]:4	param_3	XREF[2]:	08184ff0(R),
				0818500b(R)
undefined4	Stack[0x10]:4	param_4	XREF[1]:	08184fdf(R)
undefined4	Stack[0x14]:4	param_5	XREF[1]:	08184ff7(R)
undefined4	Stack[0x18]:4	param_6	XREF[1]:	08184ffe(W)
undefined4	Stack[-0x4]:4	local_4	XREF[1]:	08184fc3(R)
undefined4	Stack[-0x8]:4	local_8	XREF[1]:	08184fbb(*)
	FUN_08184fa0		XREF[2]:	0818502f(c),
				log.init:08186012(c)
08184fa0	65 8b 0d	MOV	ECX,dword ptr GS:[0x0]	
	00 00 00 00			
08184fa7	8b 89 fc	MOV	ECX,dword ptr [ECX + 0xffffffffc]	
	ff ff ff			

Figure 17: eCh0raix⁽⁹⁾ log.New function couldn't be renamed.

Figure 18 shows the lines in our script that are responsible for solving this challenge.

```
func_name = getDataAt(name_address)

#Try to define function name string.
if func_name is None:
    try:
        func_name = createAsciiString(name_address)
    except:
        print "ERROR: No name"
        continue
```

Figure 18: go_func.py.

UNRECOGNIZED STRINGS

The second issue that our scripts help to solve is related to strings within Go binaries. Let's go back to the 'Hello World' examples and take a look at the defined strings within Ghidra.

In the C binary⁽³⁾ 70 strings are defined, among which 'Hello, World!' can be found. Meanwhile, the Go binary⁽⁵⁾ includes 6,544 strings but searching for 'Hello' gives no result. Having such a high number of strings already makes it hard for reverse engineers to find the relevant ones, but in this case, the string that we would expect to find is not even recognized by Ghidra.

Location	String Value	String Representat...	Data Type
.strtab::00000000	__GNU_EH_FRAME_HDR	"__GNU_EH_FRAME_...	ds
.strtab::000000db	__GLOBAL_OFFSET_TABLE__	"__GLOBAL_OFFSET_...	ds
.strtab::00000104	__libc_csu_fini	"__libc_csu_fini"	ds
.strtab::00000114	__ITM_deregisterTMCloneTable	"__ITM_deregisterTM...	ds
.strtab::00000130	puts@@GLIBC_2.2.5	"puts@@GLIBC_2.2...	ds
.strtab::00000142	__edata	"__edata"	ds
.strtab::00000149	__libc_start_main@@GLIBC_2.2.5	"__libc_start_main...	ds
.strtab::00000168	__data_start	"__data_start"	ds
.strtab::00000175	__gmon_start__	"__gmon_start__"	ds
.strtab::00000184	__dso_handle	"__dso_handle"	ds
.strtab::00000191	__IO_stdin_used	"__IO_stdin_used"	ds
.strtab::000001a0	__libc_csu_init	"__libc_csu_init"	ds
.strtab::000001b0	__bss_start	"__bss_start"	ds
.strtab::000001bc	main	"main"	ds
.strtab::000001c1	__TMC_END__	"__TMC_END__"	ds
.strtab::000001cd	__ITM_registerTMCloneTable	"__ITM_registerTMCl...	ds
.strtab::000001e7	__cxa_finalize@@GLIBC_2.2.5	"__cxa_finalize@@G...	ds
00100001	ELF	"ELF"	ds
00100318	/lib64/ld-linux-x86-64.so.2	"/lib64/ld-linux-x86-...	ds
00100471	libc.so.6	"libc.so.6"	ds
0010047b	puts	"puts"	ds
00100480	__cxa_finalize	"__cxa_finalize"	ds
0010048f	__libc_start_main	"__libc_start_main"	ds
001004a1	GLIBC_2.2.5	"GLIBC_2.2.5"	ds
001004ad	__ITM_deregisterTMCloneTable	"__ITM_deregisterTM...	ds
001004c9	__gmon_start__	"__gmon_start__"	ds
001004d8	__ITM_registerTMCloneTable	"__ITM_registerTMCl...	ds
00102004	Hello, World!	"Hello, World!"	ds
00102061	zR	"zR"	ds

Figure 19: `world_c`⁽³⁾ defined strings with 'Hello, World!'.

Location	String Value	String Representati...	Data Type
Filter: Hello			

Figure 20: `world_go`⁽⁵⁾ defined strings without 'Hello'.

To understand the problem here, the first step is to understand what a string in Go is. Unlike in C-like languages, where strings are sequences of characters terminated with a null character, in Go strings are sequences of bytes with a fixed length. Strings are Go-specific structures, built up by a pointer to the location of the string and an integer, which is the length of the string.

```
type stringStruct struct {
    str unsafe.Pointer
    len int
}
```

Figure 21: A Go string.

These strings are stored within Go binaries as a large string blob, which consists of the concatenation of the strings without null character between them. So, while searching for 'Hello' using strings and grep gives the expected result in C, in the case of Go a huge string blob is returned containing somewhere 'Hello'.

```
> strings world_c | grep Hello
Hello, World!
```

Figure 22: world_c⁽³⁾ string search for 'Hello'.

```
> strings world_go | grep Hello
entersyscallgcBitsArenasgpcacetracehost is downillegal seekinvalid slotlfstack.pushmadvdontneedmheapSpecialmspanSpecialnot pollableraceF
inilockreleasep: m=runtime: gp=runtime: sp=short bufferspanSetSpinesweepWaiterstraceStringsuname failedwirep: p->m= != sweepgen MB) work
ers= called from failed with flushedWork heap_marked= idlethreads= is nil, not nStackRoots= s.spanclass= span.base()= syscalltick= wo
rk.nproc= work.nwait= , gp->status=, not pointer-byte block (3814697265625GC sweep waitGunjala_GondiHello, World!Masaram_GondiMende_Kika
kuiOld_HungarianSIGKILL: killSIGQUIT: quitbad flushGen bad map statedebugCall2048exchange fullfatal error: level 3 resetload64 failedmin
too largenil stackbaseout of memorysrmount errortimer expiredtraceStackTabtriggerRatio=value method xadd64 failedxchg64 failed}
```

Figure 23: world_go_println⁽¹³⁾ string search for 'Hello'.

Since the definition of strings is different, and as a result referencing them within the assembly code is also different from the usual C-like solutions, Ghidra has a hard time defining the strings within Go binaries.

The string structure can be allocated in many different ways, it can be created statically or dynamically during runtime, it varies over architecture and, even within the same architecture, multiple solutions are possible. Our team created two scripts to help Ghidra identify strings.

Dynamically allocated string structures

In the first case string structures are created at runtime. A sequence of assembly instructions is responsible for setting up the structure before a string operation. Thanks to the different instruction sets it varies across architectures. In the next few paragraphs we will go through a couple of use cases and show the instruction sequences that our script (find_dynamic_strings.py) [29] is looking for.

x86

First let's start with the 'Hello World' example⁽⁵⁾.

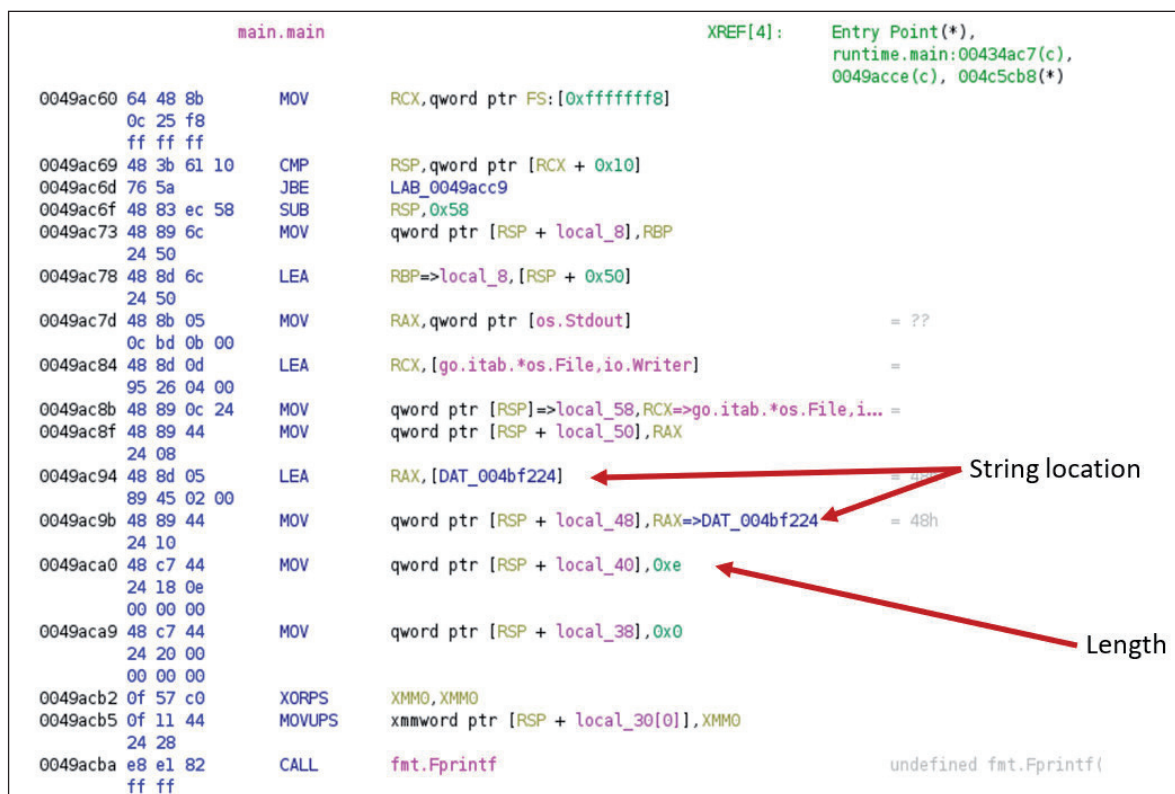


Figure 24: world_go⁽⁵⁾ dynamic allocation of string structure.

DAT_004bf224				XREF[2]:	main.main:0049ac94(*), main.main:0049ac9b(*)
004bf224	48	??	48h	H	
004bf225	65	??	65h	e	
004bf226	6c	??	6Ch	l	
004bf227	6c	??	6Ch	l	
004bf228	6f	??	6Fh	o	
004bf229	2c	??	2Ch	,	
004bf22a	20	??	20h		
004bf22b	57	??	57h	W	
004bf22c	6f	??	6Fh	o	
004bf22d	72	??	72h	r	
004bf22e	6c	??	6Ch	l	
004bf22f	64	??	64h	d	
004bf230	21	??	21h	!	
004bf231	0a	??	0Ah		

Figure 25: world_go⁽⁵⁾ undefined 'Hello, World!' string.

Figure 26 shows how the code looks after executing the script.

main.main				XREF[4]:	Entry Point(*), runtime.main:00434ac7(c), 0049acce(c), 004c5cb8(*)
0049ac60	64 48 8b	MOV	RCX,qword ptr FS:[0xffffffff8]		
	0c 25 f8				
	ff ff ff				
0049ac69	48 3b 61 10	CMP	RSP,qword ptr [RCX + 0x10]		
0049ac6d	76 5a	JBE	LAB_0049acc9		
0049ac6f	48 83 ec 58	SUB	RSP,0x58		
0049ac73	48 89 6c	MOV	qword ptr [RSP + local_8],RBP		
	24 50				
0049ac78	48 8d 6c	LEA	RBP=>local_8,[RSP + 0x50]		
	24 50				
0049ac7d	48 8b 05	MOV	RAX,qword ptr [os.Stdout]	= ??	
	0c bd 0b 00				
0049ac84	48 8d 0d	LEA	RCX,[go.itab.*os.File,io.Writer]	=	
	95 26 04 00				
0049ac8b	48 89 0c 24	MOV	qword ptr [RSP]=>local_58,RCX=>go.itab.*os.File,i...	=	
0049ac8f	48 89 44	MOV	qword ptr [RSP + local_50],RAX		
	24 08				
0049ac94	48 8d 05	LEA	RAX,[s_Hello,_World!_004bf224]	= "Hello, World!\n"	
	89 45 02 00				
0049ac9b	48 89 44	MOV	qword ptr [RSP + local_48],RAX=>s_Hello,_World!_0...	= "Hello, World!\n"	
	24 10				
0049aca0	48 c7 44	MOV	qword ptr [RSP + local_40],0xe		
	24 18 0e				
	00 00 00				
0049aca9	48 c7 44	MOV	qword ptr [RSP + local_38],0x0		
	24 20 00				
	00 00 00				
0049acb2	0f 57 c0	XORPS	XMM0,XMM0		
0049acb5	0f 11 44	MOVUPS	xmmword ptr [RSP + local_30[0]],XMM0		
	24 28				
0049acba	e8 e1 82	CALL	fmt.Fprintf	undefined fmt.Fprintf	
	ff ff				

Figure 26: world_go⁽⁵⁾ dynamic allocation of string structure after executing find_dynamic_strings.py.

The string is defined as shown in Figure 27.

s_Hello,_World!_004bf224				XREF[2]:	main.main:0049ac94(*), main.main:0049ac9b(*)
004bf224	48 65 6c	ds	"Hello, World!\n"		
	6c 6f 2c				
	20 57 6f ...				

Figure 27: world_go⁽⁵⁾ defined 'Hello, World!' string.

And 'Hello' can be found in the defined strings view in Ghidra, as shown in Figure 28.

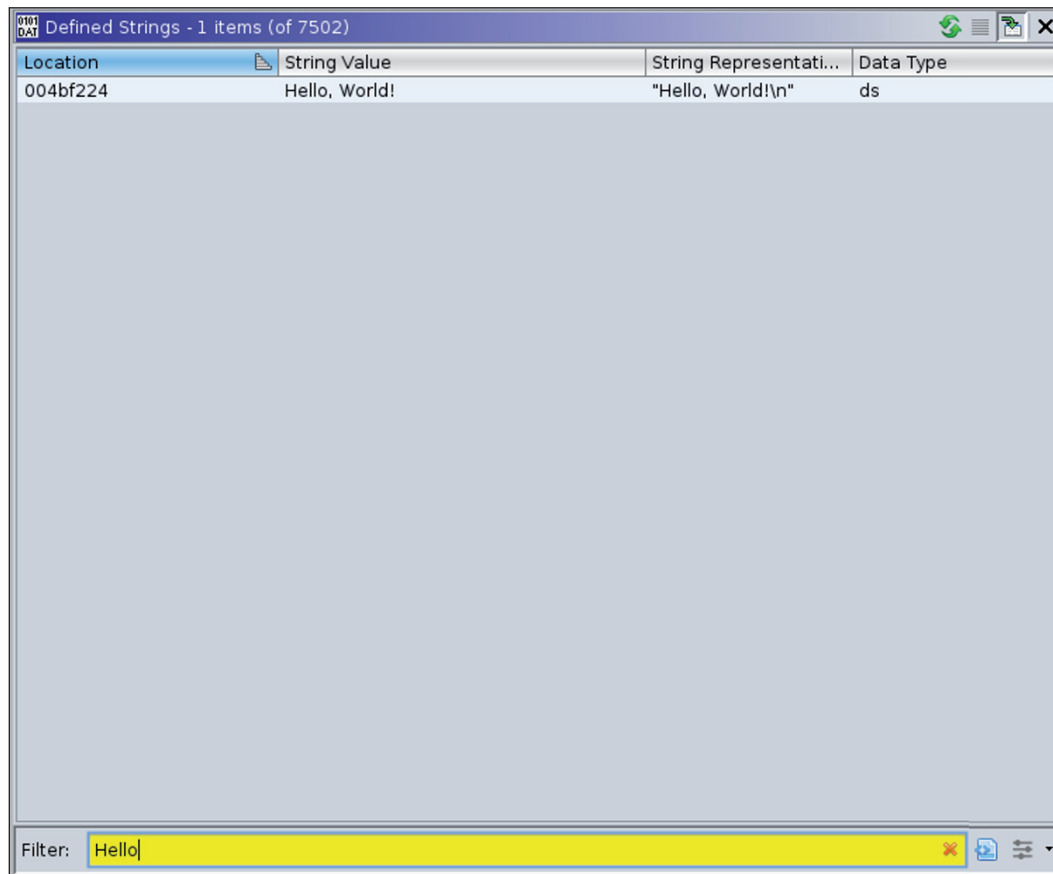


Figure 28: *world_go*⁽⁵⁾ defined strings with 'Hello'.

The script is looking for the following instruction sequences in case of 32-bit and 64-bit x86 binaries:

```
#x86
#LEA REG, [STRING_ADDRESS]
#MOV [ESP + ..], REG
#MOV [ESP + ..], STRING_SIZE
```

FUN_08208bb0		XREF[2]: 08208c3b(c), FUN_08208c40:08208cda(c)	
08208bb0	65 8b 0d 00 00 00 00	MOV	ECX,dword ptr GS:[0x0]
08208bb7	8b 89 fc ff ff ff	MOV	ECX,dword ptr [ECX + 0xffffffffc]
08208bbd	3b 61 08	CMP	ESP,dword ptr [ECX + 0x8]
08208bc0	76 74	JBE	LAB_08208c36
08208bc2	83 ec 1c	SUB	ESP,0x1c
08208bc5	c7 04 24 00 00 00 00	MOV	dword ptr [ESP]=>local_1c,0x0
08208bcc	8b 44 24 20	MOV	EAX,dword ptr [ESP + param_1]
08208bd0	89 44 24 04	MOV	dword ptr [ESP + local_18],EAX
08208bd4	8b 44 24 24	MOV	EAX,dword ptr [ESP + param_2]
08208bd8	89 44 24 08	MOV	dword ptr [ESP + local_14],EAX
08208bdc	8d 05 0e de 27 08	LEA	EAX,[DAT_0827de0e]
08208be2	89 44 24 0c	MOV	dword ptr [ESP + local_10],EAX=>DAT_0827de0e
08208be6	c7 44 24 10 17 00 00 00	MOV	dword ptr [ESP + local_c],0x17
08208bee	e8 dd c1 e7 ff	CALL	FUN_08084dd0
08208bf3	8b 44 24 14	MOV	EAX,dword ptr [ESP + local_8]
08208bf7	8b 4c 24 18	MOV	ECX,dword ptr [ESP + local_4]

Figure 29: *eCh0raix*⁽⁹⁾ dynamic allocation of string structure.

<pre>#x86_64 #LEA REG, [STRING_ADDRESS] #MOV [RSP + ..], REG #MOV [RSP + ..], STRING_SIZE</pre>					
0049ac94	48 8d 05	LEA	RAX, [DAT_004bf224]	= 48h	H
	89 45 02 00				
0049ac9b	48 89 44	MOV	qword ptr [RSP + local_48], RAX=>DAT_004bf224	= 48h	H
	24 10				
0049aca0	48 c7 44	MOV	qword ptr [RSP + local_40], 0xe		
	24 18 0e				
	00 00 00				

 Figure 30: world_go⁽⁵⁾ dynamic allocation of string structure.

ARM

For the 32-bit ARM architecture an eCh0raix ransomware sample⁽¹⁰⁾ will be used to illustrate the string recovery.

001e35bc	68 23 9f e5	ldr	r2, [PTR_DAT_001e392c]	← Pointer to address containing the string location
001e35c0	10 20 8d e5	str	r2=>DAT_0025f560, [sp, #local_90]	
001e35c4	44 20 a0 e3	mov	r2, #0x44	← String location
001e35c8	14 20 8d e5	str	r2, [sp, #local_8c]	
001e35cc	18 00 8d e5	str	r0, [sp, #local_88]	
001e35d0	1c 10 8d e5	str	r1, [sp, #local_84]	
001e35d4	44 cc f9 eb	bl	runtime.concatstring3	← Length

 Figure 31: eCh0raix⁽¹⁰⁾ dynamic allocation of string structure.

PTR_DAT_001e392c		XREF[1]:		main.main:001e35bc(R)
001e392c	60 f5 25 00	addr	DAT_0025f560	

 Figure 32: eCh0raix⁽¹⁰⁾ pointer to string address.

DAT_0025f560		XREF[2]:		main.main:001e35c0(*), 001e392c(*)
0025f560	0d	??	0Dh	
0025f561	0a	??	0Ah	
0025f562	0d	??	0Dh	
0025f563	0a	??	0Ah	
0025f564	44	??	44h	D
0025f565	6f	??	6Fh	o
0025f566	20	??	20h	
0025f567	4e	??	4Eh	N
0025f568	4f	??	4Fh	O
0025f569	54	??	54h	T
0025f56a	20	??	20h	
0025f56b	72	??	72h	r
0025f56c	65	??	65h	e
0025f56d	6d	??	6Dh	m
0025f56e	6f	??	6Fh	o
0025f56f	76	??	76h	v
0025f570	65	??	65h	e
0025f571	20	??	20h	
0025f572	74	??	74h	t
0025f573	68	??	68h	h
0025f574	69	??	69h	i
0025f575	73	??	73h	s

 Figure 33: eCh0raix⁽¹⁰⁾ undefined string.

Figure 34 shows how the code looks after executing the script.

```

001e35bc 68 23 9f e5    ldr      r2,[PTR_s__Do_NOT_remove_this_file_and_NOT_001e392c]
001e35c0 10 20 8d e5    str      r2=>s__Do_NOT_remove_this_file_and_NOT_0025f560,[sp,#local_90]
001e35c4 44 20 a0 e3    mov      r2,#0x44
001e35c8 14 20 8d e5    str      r2,[sp,#local_8c]
001e35cc 18 00 8d e5    str      r0,[sp,#local_88]
001e35d0 1c 10 8d e5    str      r1,[sp,#local_84]
001e35d4 44 cc f9 eb    bl       runtime.concatstring3

```

Figure 34: *eCh0raix*⁽¹⁰⁾ dynamic allocation of string structure after executing *find_dynamic_strings.py*.

The pointer is renamed, and the string is defined:

```

PTR_s__Do_NOT_remove_this_file_and_NOT_001e392c XREF[1]:    main.main:001e35bc(R)
001e392c 60 f5 25 00    addr      s__Do_NOT_remove_this_file_and_NOT_0025f560

```

Figure 35: *eCh0raix*⁽¹⁰⁾ pointer to string address after executing *find_dynamic_strings.py*.

```

s__Do_NOT_remove_this_file_and_NOT_0025f560 XREF[2]:    main.main:001e35c0(*),
0025f560 0d 0a 0d      ds          "\r\n\r\nDo NOT remove this file and NOT remove last line in this file!\r\n"
          0a 44 6f
          20 4e 4f ...

```

Figure 36: *eCh0raix*⁽¹⁰⁾ defined string after executing *find_dynamic_strings.py*.

The script is looking for the following instruction sequence in case of 32-bit ARM binaries:

```

#ARM, 32-bit
#LDR REG, [STRING_ADDRESS_POINTER]
#STR REG, [SP, ..]
#MOV REG, STRING_SIZE
#STR REG, [SP, ..]

```

Figure 37: The instruction sequence the script looks for.

For the 64-bit ARM architecture a Kaiji sample⁽¹²⁾ will be used to illustrate the string recovery. Here, two instruction sequences are used that only differ in one instruction.

```

LAB_0020b59c
0020b59c 00 04 00 b0    adrp     x0,0x28c000
0020b5a0 00 c4 1c 91    add      x0,x0,#0x731
0020b5a4 e0 07 00 f9    str      x0=>DAT_0028c731,[sp,#local_68]
0020b5a8 e0 07 7e b2    orr      x0,xzr,#0xc
0020b5ac e0 0b 00 f9    str      x0,[sp,#local_60]
0020b5b0 e4 d3 ff 97    bl       ddos.PathExists

LAB_0020b5bc
0020b5bc 00 04 00 f0    adrp     x0,0x28e000
0020b5c0 00 84 28 91    add      x0,x0,#0xa21
0020b5c4 e0 07 00 f9    str      x0=>DAT_0028ea21,[sp,#local_68]
0020b5c8 80 02 80 d2    mov      x0,#0x14
0020b5cc e0 0b 00 f9    str      x0,[sp,#local_60]
0020b5d0 dc d3 ff 97    bl       ddos.PathExists
0020b5d4 e0 63 40 39    ldrb     w0,[sp,#local_58]
0020b5d8 80 00 00 b5    cbnz     x0,LAB_0020b5e8

```

XREF[2]: 0020b814(j), 0020b988(j)

String location

Length

XREF[2]: 0020b680(j), 0020b7f4(j)

String location

Length

Figure 38: *Kaiji*⁽¹²⁾ dynamic allocation of string structure.

Figure 39 shows how the code looks after executing the script.

		LAB_0020b59c	XREF[2]:	0020b814(j), 0020b988(j)
0020b59c	00 04 00 b0	adrp	x0, 0x28c000	
0020b5a0	00 c4 1c 91	add	x0, x0, #0x731	
0020b5a4	e0 07 00 f9	str	x0=>s_/etc/init.d/_0028c731, [sp, #local_68]	
0020b5a8	e0 07 7e b2	orr	x0, xzr, #0xc	
0020b5ac	e0 0b 00 f9	str	x0, [sp, #local_60]	
0020b5b0	e4 d3 ff 97	bl	ddos.PathExists	
0020b5b4	e0 63 40 39	ldrb	w0, [sp, #local_58]	
0020b5b8	60 05 00 b5	cbnz	x0, LAB_0020b664	
		LAB_0020b5bc	XREF[2]:	0020b680(j), 0020b7f4(j)
0020b5bc	00 04 00 f0	adrp	x0, 0x28e000	
0020b5c0	00 84 28 91	add	x0, x0, #0xa21	
0020b5c4	e0 07 00 f9	str	x0=>s_/etc/systemd/system/_0028ea21, [sp, #local_68]	
0020b5c8	80 02 80 d2	mov	x0, #0x14	
0020b5cc	e0 0b 00 f9	str	x0, [sp, #local_60]	
0020b5d0	dc d3 ff 97	bl	ddos.PathExists	
0020b5d4	e0 63 40 39	ldrb	w0, [sp, #local_58]	
0020b5d8	80 00 00 b5	cbnz	x0, LAB_0020b5e8	

Figure 39: Kaiji⁽¹²⁾ dynamic allocation of string structure after executing `find_dynamic_strings.py`.

The strings are defined:

		s_/etc/init.d/_0028c731	XREF[1]:	main.runkshe11:0020b5a4(*)
0028c731	2f 65 74	ds	"/etc/init.d/"	
	63 2f 69			
	6e 69 74 ...			
		s_/etc/systemd/system/_0028ea21	XREF[1]:	main.runkshe11:0020b5c4(*)
0028ea21	2f 65 74	ds	"/etc/systemd/system/"	
	63 2f 73			
	79 73 74 ...			

Figure 40: Kaiji⁽¹²⁾ defined strings after executing `find_dynamic_strings.py`.

The script is looking for the following instruction sequences in case of 64-bit ARM binaries:

```
#ARM, 64-bit - version 1
#ADRP REG, [STRING_ADDRESS_START]
#ADD REG, REG, INT
#STR REG, [SP, ..]
#ORR REG, REG, STRING_SIZE
#STR REG, [SP, ..]

#ARM, 64-bit - version 2
#ADRP REG, [STRING_ADDRESS_START]
#ADD REG, REG, INT
#STR REG, [SP, ..]
#MOV REG, STRING_SIZE
#STR REG, [SP, ..]
```

Figure 41: The instruction sequence the script looks for.

As the above examples show, after executing the script, dynamically allocated string structures can be recovered. This gives a great help to reverse engineers trying to read the assembly code or look for interesting strings within the defined string window in Ghidra.

Challenges

The biggest drawback of this approach is that for each architecture, and even for different solutions within the same architecture, a new branch has to be added to the script. Also, it is very easy to evade these predefined instruction sets. In the example shown in Figure 42, in a Kaiji 64-bit ARM malware sample⁽¹²⁾ the length of the string is moved to a register earlier than our script would expect, therefore this string will be missed.

001fd734	21 01 80 d2	mov	param_2,#0x9	← Length
001fd738	e1 4b 00 f9	str	param_2,[sp, #local_c0]	
001fd73c	62 04 00 d0	adrp	param_3,0x28b000	
001fd740	42 fc 2f 91	add	param_3=>DAT_0028bbff,param_3,#0xbff	
001fd744	e2 4f 00 f9	str	param_3=>DAT_0028bbff,[sp, #local_b8]	
001fd748	e1 53 00 f9	str	param_2,[sp, #local_b0]	← String location

Figure 42: Kaiji⁽¹²⁾ dynamic allocation of string structure in an unusual way.

DAT_0028bbff				XREF[6]:	ddos.sshgo:001fd740(*), ddos.sshgo:001fd744(*), ddos.sshgo:001fd788(*), ddos.sshgo:001fd7a4(*), ddos.sshgo:001fd7c0(*), ddos.sshgo:001fd7dc(*)
0028bbff	6c	??	6Ch	l	
0028bc00	69	??	69h	i	
0028bc01	6e	??	6Eh	n	
0028bc02	75	??	75h	u	
0028bc03	78	??	78h	x	
0028bc04	5f	??	5Fh	_	
0028bc05	61	??	61h	a	
0028bc06	72	??	72h	r	
0028bc07	6d	??	6Dh	m	

Figure 43: Kaiji⁽¹²⁾ undefined string.

Statically allocated string structures

In the next case our script (find_static_strings.py) [30] looks for string structures that are statically allocated, meaning the string pointer is followed by the string length within the data section of the code.

To illustrate this let's look at the x86 eCh0raix ransomware sample⁽⁹⁾.

PTR_DAT_08436680				XREF[2]:	0820a330(*), 08431db0(*)
08436680	e1 85 27 08	addr	DAT_082785e1		
08436684	04	??	04h	← String pointers	
08436685	00	??	00h		
08436686	00	??	00h		
08436687	00	??	00h		
08436688	9d	??	9Dh		
08436689	84	??	84h	← String pointers	
0843668a	27	??	27h		
0843668b	08	??	08h		
0843668c	04	??	04h		
0843668d	00	??	00h		
0843668e	00	??	00h		
0843668f	00	??	00h		
08436690	b1	??	B1h		
08436691	84	??	84h	← String length	
08436692	27	??	27h		
08436693	08	??	08h		
08436694	04	??	04h		
08436695	00	??	00h		
08436696	00	??	00h		
08436697	00	??	00h		

Figure 44: eCh0raix⁽⁹⁾ static allocation of string structures.

In Figure 44 string pointers are followed by string length values, however Ghidra couldn't recognize the addresses or the integer data types, with the exception of the first pointer, which is directly referenced from the code.

0820a30f	8b 44 24 20	MOV	EAX,dword ptr [ESP + 0x20]
0820a313	89 04 24	MOV	dword ptr [ESP],EAX
0820a316	8b 44 24 1c	MOV	EAX,dword ptr [ESP + 0x1c]
0820a31a	89 44 24 04	MOV	dword ptr [ESP + 0x4],EAX
0820a31e	8b 05 b0	MOV	EAX,dword ptr [PTR_PTR_DAT_08431db0]
	1d 43 08		
0820a324	8b 0d b4	MOV	ECX,dword ptr [DAT_08431db4]
	1d 43 08		
0820a32a	8b 15 b8	MOV	EDX,dword ptr [DAT_08431db8]
	1d 43 08		
0820a330	89 44 24 08	MOV	dword ptr [ESP + 0x8],EAX=>PTR_DAT_08436680
0820a334	89 4c 24 0c	MOV	dword ptr [ESP + 0xc],ECX
0820a338	89 54 24 10	MOV	dword ptr [ESP + 0x10],EDX
0820a33c	e8 df f0	CALL	FUN_08209420
	ff ff		

Figure 45: eCh0raix⁽⁹⁾ pointer.

Following the string addresses, the undefined strings can be found.

	DAT_082785e1		XREF[1]:	08436680(*)
082785e1	2e	??	2Eh	.
082785e2	64	??	64h	d
082785e3	61	??	61h	a
082785e4	74	??	74h	t
082785e5	2e	??	2Eh	.
082785e6	64	??	64h	d
082785e7	62	??	62h	b
082785e8	30	??	30h	0
082785e9	2e	??	2Eh	.
082785ea	64	??	64h	d
082785eb	62	??	62h	b
082785ec	61	??	61h	a
082785ed	2e	??	2Eh	.
082785ee	64	??	64h	d
082785ef	62	??	62h	b
082785f0	66	??	66h	f
082785f1	2e	??	2Eh	.
082785f2	64	??	64h	d
082785f3	62	??	62h	b
082785f4	6d	??	6Dh	m
082785f5	2e	??	2Eh	.
082785f6	64	??	64h	d
082785f7	62	??	62h	b
082785f8	78	??	78h	x

Figure 46: eCh0raix⁽⁹⁾ undefined strings.

After executing the script, the string addresses will be defined, along with the string length values and the strings themselves.

	PTR_s_.dat_08436680		XREF[2]:	0820a330(*), 08431db0(*)
08436680	e1 85 27 08	addr	s_.dat_082785e1	
08436684	04 00 00 00	int	4h	
08436688	9d 84 27 08	addr	s_.1st_0827849d	
0843668c	04 00 00 00	int	4h	
08436690	b1 84 27 08	addr	s_.602_082784b1	
08436694	04 00 00 00	int	4h	
08436698	e5 82 27 08	addr	s_.7z_082782e5	
0843669c	03 00 00 00	int	3h	
084366a0	17 90 27 08	addr	s_.7-zip_08279017	
084366a4	06 00 00 00	int	6h	
084366a8	c1 84 27 08	addr	s_.abw_082784c1	
084366ac	04 00 00 00	int	4h	
084366b0	c5 84 27 08	addr	s_.act_082784c5	
084366b4	04 00 00 00	int	4h	
084366b8	11 8c 27 08	addr	s_.adoc_08278c11	
084366bc	05 00 00 00	int	5h	
084366c0	d9 84 27 08	addr	s_.aim_082784d9	
084366c4	04 00 00 00	int	4h	
084366c8	e1 84 27 08	addr	s_.ans_082784e1	
084366cc	04 00 00 00	int	4h	

Figure 47: eCh0raix⁽⁹⁾ static allocation of string structures after executing find_static_strings.py.

082785e1	2e 64 61 74	s_.dat_082785e1 ds	".dat"	XREF[2]:	08436680(*), 084378b0(*)
082785e5	2e 64 62 30	s_.db0_082785e5 ds	".db0"	XREF[1]:	08437248(*)
082785e9	2e 64 62 61	s_.dba_082785e9 ds	".dba"	XREF[1]:	08437250(*)
082785ed	2e 64 62 66	s_.dbf_082785ed ds	".dbf"	XREF[1]:	08437258(*)
082785f1	2e 64 62 6d	s_.dbm_082785f1 ds	".dbm"	XREF[1]:	08436ed8(*)
082785f5	2e 64 62 78	s_.dbx_082785f5 ds	".dbx"	XREF[1]:	08437260(*)
082785f9	2e 64 63 72	s_.dcr_082785f9 ds	".dcr"	XREF[2]:	084369d8(*), 08437268(*)
082785fd	2e 64 65 72	s_.der_082785fd ds	".der"	XREF[2]:	08436d70(*), 08437270(*)

Figure 48: eCh0raix⁽⁹⁾ defined strings after executing `find_static_strings.py`.

Challenges

To eliminate false positives we limit the string length, search only for printable characters, and only in data sections of the binaries. Obviously, as a result of these limitations strings can easily be missed. If you use the script feel free to experiment with it, change the values and find the best settings for your analysis. The following lines in the code are responsible for the length and character set limitations:

```
#Look for strings with printable characters only to eliminate FPs.
def isPrintable(s, l):
    for i in range(l):
        if getByte(s) not in range(32,126):
            return False
        s = s.add(1)
    return True
```

Figure 49: `find_static_strings.py`.

```
length = getInt(length_address)
#Set the possible length to eliminate FPs.
if length not in range(1,100):
    continue
```

Figure 50: `find_static_strings.py`.

Further challenges in string recovery

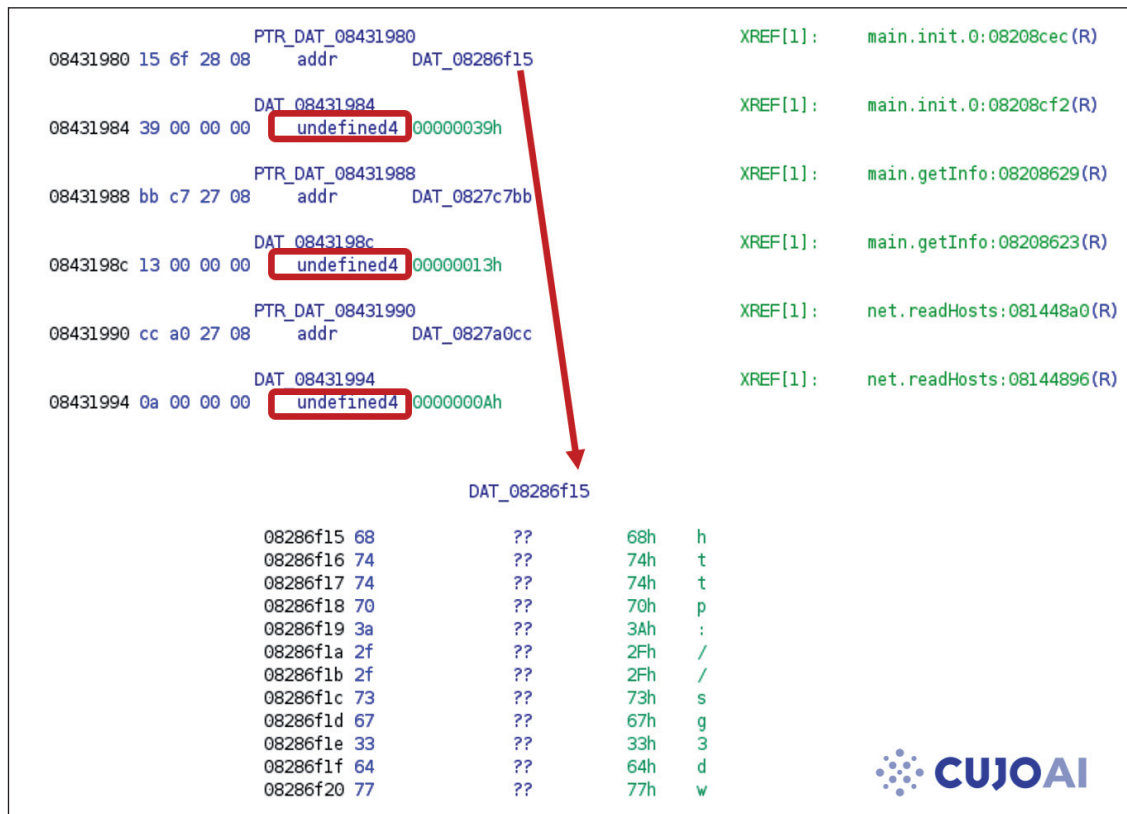
Ghidra auto analysis can falsely identify certain data types. When this happens, our script will fail to create the correct data at that specific location. To overcome this issue, first the incorrect data type has to be removed, then the new one can be created.

As an example, let's take a look at the eCh0riax ransomware⁽⁹⁾ with statically allocated string structures. Figure 51 shows the static allocation of string structures.

Here, the addresses are correctly identified, however the string length values, that are supposed to be integer data types, are falsely defined as undefined values.

Figure 52 shows the lines in our script that are responsible for removing the incorrect data types.

As shown in Figure 53, after executing the script all the data types are correctly identified and the strings are defined.

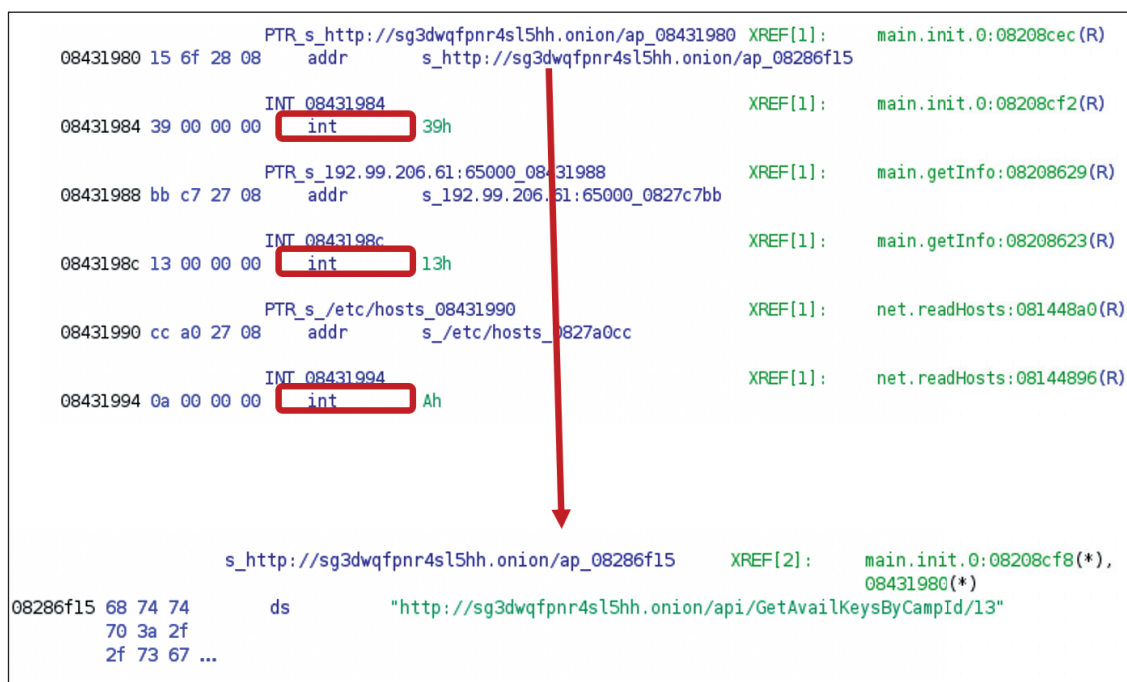
Figure 51: eCh0raix⁽⁹⁾ static allocation of string structures.

```

if getDataAt(length_address) is not None:
    data_type = getDataAt(length_address).getDataTypes()
    #Remove undefined data to be able to create int.
    #Keep an eye on other predefined data types.
    if data_type.getName() in ["undefined4", "undefined8"]:
        removeData(getDataAt(length_address))

```

Figure 52: find_static_strings.py.

Figure 53: eCh0raix⁽⁹⁾ static allocation of string structures after executing find_static_strings.py.

Another issue comes from the fact that in Go binaries strings are stored concatenated, in a large string blob. In certain cases, Ghidra define these blobs as one string. These can be identified by the high number of offcut references. Offcut references are references to certain parts of the defined string, not the address where the string starts, rather somewhere inside the string.

The example shown in Figures 54 and 55 is from an ARM Kaiji sample⁽¹²⁾.

```

s_runtime:panic_before_malloc_he_002978ff      runtime.casgstatus:00043ef4(*),
s_runtime:"*+*+####@@@!!!first path segment in URL cannot contain colon\n-s /etc/rc.d/init.d/linux_kill
s_runtime:/etc/rc.d/rcmath/big: mismatched montgomery number lengthsmemory reservation exceeds address space
s_slice:limitpanicwrap: unexpected string after type name: reflect.Value.Slice: slice index out of
s_ssh:boundsreflect: nil type passed to Type.ConvertibleToreleased less than one physical page of
s_sys:memoryruntime: debugCallV1 called by unknown caller runtime: failed to create new OS thread (have
s_tls:runtime: name offset base pointer out of rangerruntime: panic before malloc heap
s_led:initialized\nruntime: text offset base pointer out of rangerruntime: type offset base pointer out of
s_tls:rangeslice bounds out of range [%x] with length %yssh: unmarshal error for field %s of type
s_tls:%s%ssstopTheWorld: not stopped (status != _Pgcstop)syGrow bounds not aligned to palloChunkBytestls:
s_tls:failed to parse certificate from server: tls: received new session ticket from a clienttls: server
s_x509:chose an unconfigured cipher suitetls: server did not echo the legacy session IDx509: failed to
s_x509:parse rfc822Name constraint %qx509: failed to unmarshal elliptic curve pointx509: invalid elliptic
s_x509:curve private key valueP has cached GC work at end of mark terminationattempting to link in too many
s_P:shared librariesbufio: reader returned negative count from Readchacha20poly1305: message
s_at:authentication failedcurve25519: global Basepoint value was modifiedexplicit string type given to
s_buf:non-string memberfirst record does not look like a TLS handshakeslice bounds out of range [%x]
s_cha:with length %ytls: incorrect renegotiation extension contentstls: internal error: pskBinders length
s_cu:mismatchtls: server selected TLS 1.3 in a renegotiationtls: server sent two HelloRetryRequest
s_ex:messagesx509: internal error: IP SAN %x failed to parsebufio: writer returned negative count from
Writecrypto/rsa: key size too small for PSS signaturefailed to parse certificate %#d in the chain:
%wparsing/packing of this type isn't available yetruntime: cannot map pages l...
002976f3 2a 2d 2b
          2a 2d 2b
          23 23 23 ...

```

Figure 54: Kaiji⁽¹²⁾ falsely defined string in Ghidra.

```

s_runtime:panic_before_malloc_he_002978ff      runtime.casgstatus:00043ef4(*),
s_runtime:text_offset_base_pointe_0029792d      runtime.doInit:0004eefc(*),
s_runtime:type_offset_base_pointe_0029795b      runtime.sigpanic:00055da4(*),
s_slice:bounds_out_of_range [%x]_w_00297989      runtime.sigpanic:00055de4(*),
s_ssh:unmarshal_error_for_field_%_002979b7      runtime.sigpanic:00055f24(*),
s_sysGrow:bounds_not_aligned_to_pa_00297a13      runtime.sigpanic:00055f64(*),
s_tls:failed_to_parse_certificate_00297a41      runtime.getStackMap:0005a7d4(*),
s_led:to_parse_certificate_from_se_00297a49      runtime.morestackcc:0005a834(*),
s_tls:received_new_session_ticket_00297a6f      runtime.resolveNameOff:00065b1c(...
s_tls:server_chose_an_unconfigure_00297a9d
s_tls:server_did_not_echo_the_leg_00297acb
s_x509:failed_to_parse_rfc822Name_00297af9
s_x509:failed_to_unmarshal_ellipt_00297b27
s_x509:invalid_elliptic_curve_pri_00297b55
s_P:has_cached_GC_work_at_end_of_m_00297b83
s_attempting_to_link_in_too_many_s_00297bb2
s_bufio:reader_returned_negative_c_00297be1
s_chacha20poly1305:message_authen_00297c10
s_curve25519:global_Basepoint_val_00297c3f
s_explicit_string_type_given_to_no_00297c6e
002976f3 2a 2d 2b      ds      "*+*+####@@@!!!first path segment in URL cannot contain colon\n-s /etc/rc.d
          2a 2d 2b
          23 23 23 ...

```

Figure 55: Kaiji⁽¹²⁾ offcut references of a falsely defined string.

To find falsely defined strings, one can use the defined strings window of Ghidra and sort the strings by offcut reference count. Large strings with numerous offcut references can be undefined manually before executing the string recovery scripts, so the scripts can successfully create the correct string data types. Figure 56 shows Kaiji's defined strings.

Finally, we will show an issue in versions of Ghidra decompiler view prior to version 9.2. Once a string is successfully defined, either manually or by one of our scripts, it will be nicely visible in the listing view of Ghidra, giving a great help to reverse engineers when reading the assembly code. However, the decompiler view in earlier versions of Ghidra couldn't handle fixed length strings correctly and, regardless of the length of the string, it would display everything until it found a null character. Thankfully this issue was solved in Ghidra 9.2.

The issue is illustrated in Figures 57 and 58 using the eCh0raix sample⁽⁹⁾.

Defined Strings - 10814 items				
Location	String Value	Data Type	Byte Count	Offset Reference Count
0022073d	certificateAuthorities	ds	23	1
00220ec1	ReplaceAllLiteralString	ds	24	1
00220ef5	responseMessageReceived	ds	24	1
00220f29	verifyServerCertificate	ds	24	1
00221561	hashForClientCertificate	ds	25	1
00221e1e	asn1:"explicit,tag:1"	ds	22	1
00221e53	handlePostHandshakeMessage	ds	27	1
00222552	secureRenegotiationSupported	ds	30	1
00222ebd	asn1:"optional,tag:2"	ds	23	1
00290069	ckunpa	ds	6	1
002903f7	queuefinalizer during GC	ds	24	1
00330cff	runtime.dropg	ds	14	1
00460248	-----END	ds	12	1
00460258	-----BEGIN	ds	16	1
0029bb9c	0001020304050607080910111...	ds	969	2
002e9100	expand 32-byte k	ds	20	3
002e91a0	expand 32-byte k	ds	20	3
00293a08	3552713678800500929355621...	ds	170	4
0028b3b3	= is not mcount= minutes nallo...	ds	225	23
002976f3	*-+*+####@@@@@!!!!first pat...	ds	4517	95

Figure 56: Kaiji⁽¹²⁾ defined strings.

main.checkReadmeExists		XREF[2]: 08208c3b(c), main.init.0:08208cda(c)	
08208bb0	65 8b 0d MOV ECX,dword ptr GS:[0x0]		
08208bb7	8b 89 fc MOV ECX,dword ptr [ECX + 0xffffffffc]		
08208bbd	3b 61 08 CMP ESP,dword ptr [ECX + 0x8]		
08208bc0	76 74 JBE LAB_08208c36		
08208bc2	83 ec 1c SUB ESP,0x1c		
08208bc5	c7 04 24 MOV dword ptr [ESP]=>local_1c,0x0		
08208bcc	8b 44 24 20 MOV EAX,dword ptr [ESP + param_1]		
08208bd0	89 44 24 04 MOV dword ptr [ESP + local_18],EAX		
08208bd4	8b 44 24 24 MOV EAX,dword ptr [ESP + param_2]		
08208bd8	89 44 24 08 MOV dword ptr [ESP + local_14],EAX		
08208bdc	8d 05 0e LEA EAX,[s_/README_FOR_DECRYPT.txt_0827de0e]		
08208be2	89 44 24 0c MOV dword ptr [ESP + local_10],EAX=>s_/README_FOR_DECRYPT.txt_0827de0e		
08208be6	c7 44 24 10 17 00 MOV dword ptr [ESP + local_c],0x17		
08208bee	e8 dd c1 CALL runtime.concatstring2		

Figure 57: eCh0raix⁽⁹⁾ defined string in listing view.

[Decompile: main.checkReadmeExists] - (echoraix_test2)	
19	runtime.concatstring2
20	(0,param_1,param_2,
22	"/README_FOR_DECRYPT.txt/etc/apache2/mime.types/etc/pki/tls/cacert.pem23283064365386962890625<invalid reflect.Value>CPU time limit exceededLogical_Order_ExceptionMB during sweep: swept Noncharacter_Code_PointSIGIO: i/o now possibleSIGSYS: bad system callVariant Also Negotiatesacquirep: already in goasn1: structure error: bytes.Buffer: too largechan receive (nil chan)close of closed channelcommand not implementeddevice or resource busyfatal: morestack on go\inflate: internal error: garbage collection scangcDrain phase incorrecthttp2: handler panickedhttp2: invalid trailershttp: request too largeinterrupted system callinvalid URI for requestinvalid m->lockedInt = json: cannot unmarshal left over markroot jobsmakechan: bad alignmentmalformed HTTP responsemissing port in addressmissing protocol schememissing type in runfqmisuse of profBuf.writenanotime returning zeronet/http: abort Handlernetwork not implementedno application protocolno space left on devicenon-zero reserved fieldoperation not permittedoperation not supportedpanic during preemptoffprogresize: invalid argprofilng timer expiredreflect.Value.Interfacereflect.Value.NumMethodreflect.methodValueCallruntime: internal errorruntime: invalid type runtime: netpoll failedruntime: s.allocCount=s.allocCount > s.nelemsschedule: holding locksssegment length too longskipping Question Classspan has no free stacksgrowth after forksyntax error in patterntext/css; charset=utf-8text/xml; charset=utf-8time: invalid duration too many pointers (>10)truncated tag or lengthunexpected address typeunexpected signal valueunknown error code 0x%unlock of unlocked lockunpacking Question.Nameunpacking Question.Typeunsupported certificatevarint integer overflowwork.nwait > work.nproc%v.WithDeadline(%s [%s])/usr/share/lib/zoneinfo/116415321826934814453125582076609134674072265625Request Entity Too Largebad defer entry in panicbad defer size class: i=block index out of rangeCan't scan our own stackconnection reset by peerdouble traceGCSweepStartererror decrypting messagefile size li..." /* TRUNCATED STRING LITERAL */ ,0x17);

Figure 58: eCh0raix⁽⁹⁾ defined string in decompile view in Ghidra 9.1.

FUTURE WORK

In this paper we proposed solutions for two issues within Go binaries to help reverse engineers when they are using Ghidra to statically analyse malware written in Go. In the first topic we discussed how to recover function names in stripped Go binaries. Then we proposed multiple solutions for defining strings within Ghidra. The scripts that we created and files we used for the examples in this paper are publicly available, the links can be found below.

There are even more possibilities to aid Go reverse engineering – the two topics that we discussed here are just the beginning. As a next step we are planning to dive deeper into Go function call conventions and types system.

In Go binaries arguments and return values are passed to functions using the stack, rather than registers. Currently, Ghidra has a hard time correctly detecting these. Helping Ghidra to support Go's calling convention will help reverse engineers to understand the purpose of the analysed functions.

The other interesting topic is types within Go binaries. Just as it was possible to extract function names from the investigated files, Go binaries also store information about the used types. Recovering these types can be a great help during reverse engineering. In the example shown in Figures 59 – 61 we recovered the main.Info structure in an eCh0raix ransomware sample⁽⁹⁾. This structure tells us what information the malware is expecting from the C2 server.

main.info_struct			XREF[3]:	main.getInfo:082085fc(*), main.getInfo:08208602(*), 08225100(*)
0824bd20	10 00 00 00	ddw	10h	
0824bd24	0c 00 00 00	ddw	Ch	
0824bd28	15 e7 c0 27	ddw	27C0E715h	
0824bd2c	07	db	7h	
0824bd2d	04	db	4h	
0824bd2e	04	db	4h	
0824bd2f	19	db	19h	
0824bd30	28 c8 20 08	addr	PTR_PTR_type..hash.main.Info_0820c828	
0824bd34	fc a0 2b 08	addr	DAT_082ba0fc	
0824bd38	20 75 00 00	ddw	7520h	
0824bd3c	e0 a0 01 00	ddw	1A0E0h	
0824bd40	00 00 00 00	ddw	0h	
0824bd44	60 bd 24 08	addr	PTR_rsapublickey_structfield_0824bd60	
0824bd48	02 00 00 00	ddw	2h	
0824bd4c	02 00 00 00	ddw	2h	
0824bd50	5c 0d 00 00	ddw	D5Ch	
0824bd54	00 00	dw	0h	
0824bd56	00 00	dw	0h	
0824bd58	28 00 00 00	ddw	28h	
0824bd5c	00 00 00 00	ddw	0h	

Figure 59: eCh0raix⁽⁹⁾ main.info structure.

PTR_rsapublickey_structfield_0824bd60			XREF[1]:	0824bd44(*)
0824bd60	60 aa 22 08	addr	rsapublickey_structfield	
0824bd64	a0 a7 23 08	addr	string_type	
0824bd68	00 00 00 00	ddw	0h	
0824bd6c	18 cf 21 08	addr	readme_structfield	
0824bd70	a0 a7 23 08	addr	string_type	
0824bd74	10 00 00 00	ddw	10h	

Figure 60: eCh0raix⁽⁹⁾ main.info fields.

```
type main.Info struct{
    RsaPublicKey string
    Readme string
}
```

Figure 61: eCh0raix⁽⁹⁾ main.info structure.

As these examples illustrated there are still a lot of interesting areas to discover within Go binaries from reverse engineering point of view.

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GITHUB REPOSITORY WITH SCRIPTS AND ADDITIONAL MATERIALS

- <https://github.com/getCUJO/ThreatIntel/tree/master/Scripts/Ghidra>
- https://github.com/getCUJO/ThreatIntel/tree/master/Research_materials/Golang_reversing

FILES USED DURING THE RESEARCH

	File name	SHA-256
(1)	world.c	761301bb14ea3b678650fc1b6da768f009387ee726712e291d57e2d7985613d0
(2)	world.go	7cb3316a7b89eb996e8dbb0d0fb277136cd588cc54642f3b09aa84cd177cb3a2
(3)	world_c	76a5c4ef9277b97660f2c412e67ff2c3826e699913db86cd333e8f1d4fb5b8a3
(4)	world_c_strip	486a93362a6a8bc3b449fd6ba07656011c687ed31a19091c329a434bfff4d75bb
(5)	world_go	d0d4781de4ffd5f8e18d59328eccd373a782eecd55a2c5199b7dc6598cfb99e
(6)	world_go_strip	9b975bd9406a8b79a414195e184be0c82bb1593979577f0344c797f9bcd4ad0b
(7)	world_go.exe	9e36291f5fc67fdb9e5e17b636d34b39f2cc39f328916a9012a8f8d545e9d0c8
(8)	world_go_strip.exe	c5b66623942a0cea6df30541e92afe93172be7bb4dbdd42a1fa354e9edd79a1d
(9)	eCh0raix - x86	154dea7cace3d58c0cecccb5a3b8d7e0347674a0e76daffa9fa53578c036d9357
(10)	eCh0raix - ARM	3d7ebe73319a3435293838296fbb86c2e920fd0ccc9169285cc2c4d7fa3f120d
(11)	Kaiji - x86_64	f4a64ab3fffc0b4a94fd07a55565f24915b7a1aaec58454df5e47d8f8a2eec22a
(12)	Kaiji - ARM	3e68118ad46b9eb64063b259fca5f6682c5c2cb18fd9a4e7d97969226b2e6fb4
(13)	world_go_println	fa00f5ad2aa79a6245a28516bc285ae8c36f075d818787aadff6f3e850e2ec5c

SOLUTIONS BY OTHER RESEARCHERS FOR VARIOUS TOOLS

IDA Pro

- <https://github.com/sibears/IDAGolangHelper>
- https://github.com/strazzere/golang_loader_assist

radare2 / Cutter

- <https://github.com/f0rki/r2-go-helpers>
- https://github.com/JacobPimental/r2-gohelper/blob/master/golang_helper.py
- <https://github.com/CarveSystems/gostringsr2>

Binary Ninja

- <https://github.com/f0rki/bn-goloader>

Ghidra

- <https://github.com/felberj/gotools>
- https://github.com/ghidraninja/ghidra_scripts/blob/master/golang_renamer.py